

Memory management

Pages - fixed-size block of logical memory

Frames - fixed-size block of physical memory

page table - to translate logical to physical addresses

Hit ratio(α) – percentage of times that a page number is found in the associative registers

- Consider $\alpha = 80\%$, 20ns for TLB search, 100ns for memory access
- $EAT = 0.80 \times (20+100) + 0.20 \times (20+100+100) = 140\text{ns}$
- Consider hit ratio $\rightarrow \alpha = 98\%$, 20ns for TLB search, 100ns for memory access
- $EAT = 0.98 \times 120 + 0.02 \times 220 = 122\text{ns}$

Q. The memory access time is 1 nanosecond for a read operation with a hit in cache, 5 nanoseconds for a read operation with a miss in cache, 2 nanoseconds for a write operation with a hit in cache and 10 nanoseconds for a write operation with a miss in cache. Execution of a sequence of instructions involves 100 instruction fetch operations, 60 memory operand read operations and 40 memory operand write operations. The cache hit-ratio is 0.9. The average memory access time (in nanoseconds) in executing the sequence of instructions is _____.

Sol. Total instruction=100 instruction fetch operation +60 memory operand read operation +40 memory operand write op

Total instruction=200 instructions (operations)

Time taken for fetching and reading instructions= $160(0.9(1)+0.1(5)) = 224\text{ns}$

Memory operands Write operations time = $40(0.9(2)+0.1(10)) = 112\text{ns}$

Total time taken for executing 200 instructions = $224 + 112 = 336\text{ns}$

Average memory access time = $336\text{ ns}/200=1.68\text{ns}$

Q. Let the page fault service time to 10 ms in a computer with average memory access time being 20 ns. If one page fault is generated for every 10^6 memory accesses, what is the effective access time for the memory?

Sol.

P = page fault rate

$$P = 1/10^6$$

$$\begin{aligned} \text{EA} &= p \times \text{page fault service time} + (1 - p) \times \text{Memory access time} \\ &= 1/10^6 \times 10 \times 10^6 + (1 - 1/10^6) \times 20 \approx 29.9 \text{ ns} \end{aligned}$$

Belady's anomaly

It is a phenomenon where increasing the number of page frames results in an increase in the number of page faults for a given memory access pattern.

- The reference String is given as 0 1 5 3 0 1 4 0 1 5 3 4.

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Case 1: Number of frames = 3

Request	0	1	5	3	0	1	4	0	1	5	3	4
Frame 3			5	5	5	1	1	1	1	1	3	3
Frame 2		1	1	1	0	0	0	0	0	5	5	5
Frame 1	0	0	0	3	3	3	4	4	4	4	4	4
Miss/Hit	Miss	Miss	Miss	Miss	Miss	Miss	Miss	Hit	Hit	Miss	Miss	Hit

Number of Page Faults = 9

- The reference String is given as 0 1 5 3 0 1 4 0 1 5 3 4.

Case 2: Number of frames = 4

Request	0	1	5	3	0	1	4	0	1	5	3	4
Frame 4				3	3	3	3	3	3	5	5	5
Frame 3			5	5	5	5	5	5	1	1	1	1
Frame 2		1	1	1	1	1	1	0	0	0	0	4
Frame 1	0	0	0	0	0	0	4	4	4	4	3	3
Miss/Hit	Miss	Miss	Miss	Miss	Hit	Hit	Miss	Miss	Miss	Miss	Miss	Miss

Number of Page Faults = 10